

Outline

- 1 Defining Computational Security
- 2 Security Models for Public Key Encryption
- 3 Security Models for Digital Signatures
- 4 Security Assumptions and Results

Bounding the Adversary's Capabilities

Perfect Security vs. Computational Security

- **Perfect Security:** The adversary is unbounded, but only a few positive results can be obtained. (E.g., perfect one-time encryption.)

Turing Machine

A model for mechanic computation, consisting of

- A finite automaton

- A storage device (“tape”)

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Computation step: $(s, wr, mv) \leftarrow f(s, c)$

$wr \in \{\text{write}_0, \text{write}_1, \text{erase}\}$

$mv \in \{\text{move_left}, \text{move_right}\}$

$c =$ contents of the current cell

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Input/Output of the Turing Machine is the content of the tape in the `init/halt` state

Turing Machine

state

s

transition func.

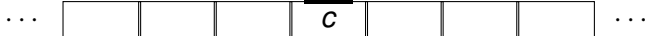
(s, c)

f

$(s^+, wr, mv) = f(s, c)$

read/write head

tape



Turing Machine

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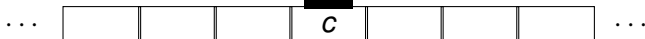
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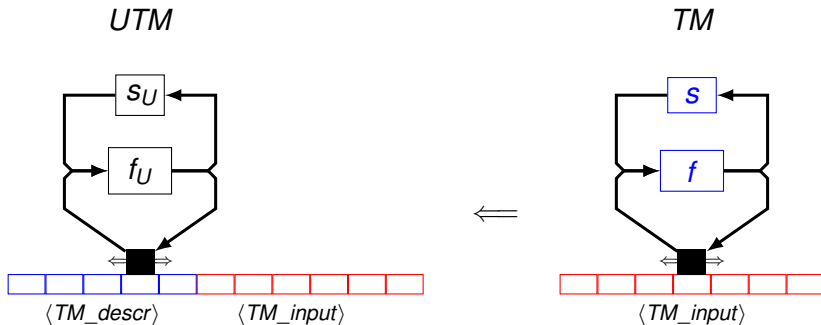


$s \in \mathcal{S} \cup \{\text{init}, \text{halt}\}$

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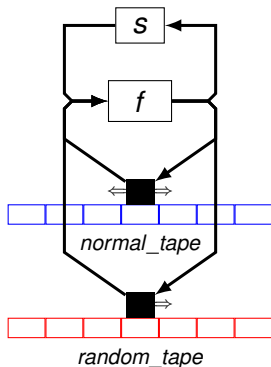
Universal Turing Machine



It can simulate any known (classical) computing device with reasonable efficiency

Probabilistic Turing Machine

PTM



A Turing Machine that take random decisions, from an additional input tape.

Computation step:

$$(s, wr, mv, mvr) \leftarrow f(s, c, r)$$

$wr \in \{\text{write}_0, \text{write}_1, \text{erase}\}$

$mv \in \{\text{move_left}, \text{move_right}\}$

$mvr \in \{\text{keep}, \text{move_right}\}$

c = contents of the current normal cell

r = contents of the current random cell

Algorithmic Complexity

A running time limitation (number of steps of the Probabilistic Turing Machine) implies similar space and randomness limitations.

Every read/write operations takes one step.

Reading a random bit takes one step.

The running time typically depends on the size of the input (description of the problem to be solved).

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Uniform approach: A single Turing Machine tries to solve problems of all sizes.

Asymptotic analysis: We study the problem complexity by the type of growth of the running time as a function of the input size.

Hardness Notions

Easy computation: For a problem family \mathcal{P} , there exists a Turing Machine TM such that

- The running time is polynomial on the size of $P \in \mathcal{P}$.
- It always outputs a correct solution of any problem instance $P \in \mathcal{P}$

$$\exists T(\cdot) \in \mathbf{poly} \quad \forall P \in \mathcal{P}$$

$$\text{time}(\text{TM}, P) \leq T(|P|) \quad \text{and} \quad \text{valid}(P, \text{TM}(P)) = 1$$

poly: The set of positive polynomials.

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A model for normal operations performed by honest parties in a protocol.

poly: The set of positive polynomials.

Hardness Notions

Feasible computation: There exists a Probabilistic Turing Machine TM such that

- The running time is polynomial on the size of $P \in \mathcal{P}$.
- The probability that it gives a correct solution for any problem instance $P \in \mathcal{P}$ is not *negligible*.

$\exists T(\cdot), Q(\cdot) \in \mathbf{poly} \quad \forall P \in \mathcal{P}$
 $\text{time}(\text{TM}, P) \leq T(|P|)$ and $\Pr(\text{valid}(P, \text{TM}(P)) = 1) \geq 1/Q(|P|)$

negl: The set of functions $f(\cdot)$ such that for any $Q(\cdot) \in \mathbf{poly}$ there exists n_0 s.t. $f(n) < 1/Q(n)$ for all $n > n_0$.

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A model for a successful attack against a protocol.

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Hardness Notions

Infeasible computation: For all Probabilistic Turing Machines TM such that the running time is polynomial on the size of $P \in \mathcal{P}$, the probability that it gives a correct solution for $P \in \mathcal{P}$ is a negligible function of its length.

If $\exists T(\cdot) \in \mathbf{poly}$ such that $\forall P \in \mathcal{P}$ $\text{time}(\text{TM}, P) \leq T(|P|)$ then $\Pr(\text{valid}(P, \text{TM}(P)) = 1) \in \mathbf{negl}(|P|)$.

PPTM: A Probabilistic Turing Machine running in Polynomial time.

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A model for a hard problem family.

PPTM: A Probabilistic Turing Machine running in Polynomial time.

Hardness Notions

Summary:

- **Easy:** There exists a TM that solves all problem instances in polynomial time.
- **Feasible:** There exists a PPTM that solves any problem instance with a non-negligible probability.
- **Infeasible:** All PPTM can only solve problem instances with a negligible probability.

In cryptography, the problem P is chosen at random among all problems with a specific size (e.g., the description of the problem instance contains a randomly generated public key).

Security as a Game

Security is modelled as a Game (random experiment) between a Challenger C and an Adversary A :

- C chooses a random problem instance of a given size
- A tries to find a solution in polynomial time
- A succeeds only with a negligible probability

Here, the success probability is computed respect to

- The internal coin tosses of A
- The random choice of C

In general, the game can have more interaction steps.

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One-Time Security for Symmetric Encryption

The adversary tries to decrypt a random ciphertext from scratch.

Challenger

$\leftarrow (\ell, \text{KeyGen}, \text{Enc}, \text{Dec}) \rightarrow$

Adversary

$k_* \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

$c_* = \text{Enc}(k_*, m_*)$; **send**(c_*)

receive(m')

accept if $m' = m_*$

receive(c_*)

???

send(m')

It is too simplistic: the adversary does not know any message / ciphertext pair for the same key. (It is a one-time model.)

One-Time Security for Symmetric Encryption

The same, written in pseudocode style (we don't know what the adversary is really doing):

Experiment $\text{Exp-SE-OT}(\Pi, \mathcal{A}, \ell)$:

$k_* \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

$c_* \leftarrow \text{Enc}(k_*, m_*)$

$m' \leftarrow \mathcal{A}(c_*)$

if $m' = m_*$ **output** 1; // \mathcal{A} wins

else output 0;

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else output 0;

Definition (SE-OT)

The symmetric encryption scheme Π is SE-OT secure if for all PPTM \mathcal{A}

$$\Pr[\text{Exp-SE-OT}(\Pi, \mathcal{A}, \ell) = 1] \in \mathbf{negl}(\ell)$$

One-Way Security for Public Key Encryption

Now, the adversary also receives the public key:
It can encrypt messages of its choice with the same target key
(CPA model)

Challenger

$\leftarrow (\ell, \text{KeyGen}, \text{Enc}, \text{Dec}) \rightarrow$

Adversary

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

$c_* = \text{Enc}(pk_*, m_*); \text{send}(pk_*, c_*)$

receive(m')

accept if $m' = m_*$

receive(pk_*, c_*)

???

send(m')

One-Way Security for Public Key Encryption

Experiment $\text{Exp-PKE-OW-CPA}(\Pi, \mathcal{A}, \ell)$:

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

$c_* \leftarrow \text{Enc}(pk_*, m_*)$

$m' \leftarrow \mathcal{A}(pk_*, c_*)$

if $m' = m_*$ **output** 1; // \mathcal{A} wins

else output 0;

Definition (PKE-OW-CPA)

The public key encryption scheme Π is PKE-OW-CPA secure if for all PPTM \mathcal{A}

$$\Pr[\text{Exp-PKE-OW-CPA}(\Pi, \mathcal{A}, \ell) = 1] \in \mathbf{negl}(\ell)$$

One-Way Security for Public Key Encryption

Most proposed PKE schemes are conjectured to be OW-CPA secure, with a proper choice of the parameters.

E.g., ElGamal, RSA, Rabin, Paillier.

Regev's PKE is not OW-CPA secure because it encrypts a single bit (thus the adversary has easily a success probability of $1/2$).

There is a known security proof showing that Rabin is PKE-OW-CPA secure if and only if factoring an RSA modulus (with $p, q = 3 \pmod{4}$) is hard.

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PKE-OW-CPA security is still unrealistic when the adversary has some a priori knowledge about the target message m_* .

Semantic Security

The adversary chooses some leakage function f for m_* :

Challenger

$\leftarrow (\ell, \text{KeyGen}, \text{Enc}, \text{Dec}) \rightarrow$

Adversary

$b \leftarrow \{0, 1\}$

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

send (pk_*)

receive $(f(\cdot))$

$m_0, m_1 \leftarrow \mathcal{M}; h = f(m_0)$

$c_* = \text{Enc}(pk_*, m_b); \text{send}(c_*, h)$

receive (b')

accept if $b' = b$

receive (pk_*)

???

send $(f(\cdot))$

(keep the internal state st)

receive (c_*, h)

???

send (b')

This is a two-stage game (the Adversary is split in two stages).

Indistinguishability of Ciphertexts

Semantic Security game is shown to be equivalent to:

Challenger

$\leftarrow (\ell, \text{KeyGen}, \text{Enc}, \text{Dec}) \rightarrow$

Adversary

$b \leftarrow \{0, 1\}$

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

send (pk_*)

receive (m_0, m_1)

$c_* = \text{Enc}(pk_*, m_b)$

send (c_*)

receive (b')

accept if $b' = b$

receive (pk_*)

???

send (m_0, m_1)

(keep the internal state st)

receive (c_*)

???

send (b')

m_0, m_1 must be valid messages. (Otherwise, the Challenger rejects).

Indistinguishability of Ciphertexts

Experiment $\text{Exp-PKE-IND-CPA}(\Pi, \mathcal{A}_1, \mathcal{A}_2, \ell)$:

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$(m_0, m_1, st) \leftarrow \mathcal{A}_1(pk_*)$ // st is the internal state passing from \mathcal{A}_1 to \mathcal{A}_2

$b \leftarrow \{0, 1\}$

$c_* \leftarrow \text{Enc}(pk_*, m_b)$

$b' \leftarrow \mathcal{A}_2(st, c_*)$

if $b' = b$ **output** 1; // $\mathcal{A} = (\mathcal{A}_1, \mathcal{A}_2)$ wins

else output 0;

Definition (PKE-IND-CPA)

The public key encryption scheme Π is PKE-IND-CPA secure if for all PPTM \mathcal{A}

$$|\Pr[\text{Exp-PKE-IND-CPA}(\Pi, \mathcal{A}, \ell) = 1] - 1/2| \in \mathbf{negl}(\ell)$$

$1/2$ is the success probability of a dummy adversary!

IND-CPA Security

Some well-known PKE schemes are conjectured to be IND-CPA secure, with a proper choice of the parameters.

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But RSA and Rabin do not achieve IND-CPA security.

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No deterministic PKE can achieve IND-CPA security.

\mathcal{A}_2 can just encrypt m_0 and compare the result with c_* .

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RSA and Rabin need some randomization (e.g., a randomized message padding scheme).

Beyond IND-CPA Security

Some realistic attacks fall outside the IND-CPA model.

The adversary has limited extra access to:

- An oracle that tells whether a (possibly manipulated) ciphertext is valid or not.
- An oracle that decrypts a (possibly manipulated) ciphertext.
- An oracle that decrypts a ciphertext related to the target one c_* .

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A maximal notion of security is defined: IND-CCA (for Chosen Ciphertext Attack).

The adversary can ask for decryptions of any possible ciphertext except for c_* .

IND-CCA Security

Challenger

 $\leftarrow (\ell, \text{KeyGen}, \text{Enc}, \text{Dec}) \rightarrow$

Adversary

 $(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$
send (pk_*)
receive (pk_*)

(many times)

receive (c_i)
 $m_i = \text{Dec}(sk_*, c_i)$; **send** (m_i)

???

send (c_i) // oracle call

receive (m_i)
receive (m_0, m_1)
 $b \leftarrow \{0, 1\}$; $c_* = \text{Enc}(pk_*, m_b)$; **send** (c_*)

???

send (m_0, m_1)
receive (c_*)

(many times)

receive (c_i)
if $c_i \neq c_*$; $m_i = \text{Dec}(sk_*, c_i)$; **send** (m_i)
else abort

???

send (c_i) // oracle call

receive (m_i)
receive (b')
accept if $b' = b$

???

send (b')

IND-CCA Security

Experiment $\text{Exp-PKE-IND-CCA}(\Pi, \mathcal{A}_1, \mathcal{A}_2, \ell)$:

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$(m_0, m_1, st) \leftarrow \mathcal{A}_1^{\mathcal{O}_1(\cdot)}(pk_*)$

$b \leftarrow \{0, 1\}$

$c_* \leftarrow \text{Enc}(pk_*, m_b)$

$b' \leftarrow \mathcal{A}_2^{\mathcal{O}_2(\cdot)}(st, c_*)$

if $b' = b$ **output** 1; // $\mathcal{A} = (\mathcal{A}_1, \mathcal{A}_2)$ wins

else output 0;

Oracle $\mathcal{O}_1(c)$ // Decryption oracle

return $\text{Dec}(sk_*, c)$

Oracle $\mathcal{O}_2(c)$ // Restricted decryption oracle

if $c \neq c_*$ **return** $\text{Dec}(sk_*, c)$ **else abort**

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The PKE Π is PKE-IND-CCA secure if for all PPTM \mathcal{A}

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- **Paillier**: The same attack also applies.
- **Hashed ElGamal**: XORing the second component of c_* with any mask z makes the oracle answer $m_b \oplus z$.
- **Regev**: Simply adding $(q - 1)/2$ to the second component of c_* makes the oracle answer $1 - m_b$.

Known CCA Attacks

With similar attacks, it is shown that:

Lemma

Any homomorphic PKE is IND-CCA insecure.

There is a tradeoff between security and functionality.

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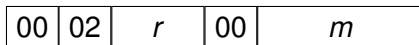
New methods are necessary to upgrade existing PKE to the IND-CCA security level.

- Randomized paddings, or Fujisaki-Okamoto Transforms (in the Random Oracle Model)
- New cryptographic tools (Hash Proof Systems, Canetti-Halevi-Katz Transform, . . .)

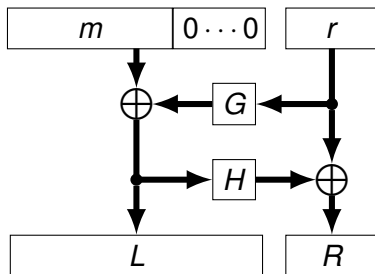
E.g., Fixing RSA PKE

- Plain RSA PKE is conjectured to be OW-CPA secure, but it is neither IND-CPA nor OW-CCA secure.
- A simple randomized message padding can give IND-CPA security, but not OW-CCA.
- Bleichenbacher's attack against PKCS#1 v1.5 shows a practical CCA attack against a padded version of RSA.
- RSA-OAEP (PKCS#1 v2) fixes the attack and provides IND-CCA security for RSA (in the Random Oracle Model).

PKCS#1 v1.5 vs. v2.0 Message Encodings



PKCS#1 v1.5



PKCS#1 v2.0

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Universal Unforgeability (UF)

The basic security notion for signatures: forge a valid signature for any given message only from the public key:

Challenger

$\leftarrow (\ell, \text{KeyGen}, \text{Sig}, \text{Ver}) \rightarrow$

Adversary

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

send (pk_*, m_*)

receive (s_*)

accept if $\text{Ver}(pk_*, m_*, s_*) = 1$

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Actually, the random selection of m_* is not well defined!.

UF-RMA Security

In UF-RMA security, the adversary can ask for valid signatures on random messages.

Experiment $\text{Exp-UF-RMA}(\Sigma, \mathcal{A}, \ell)$:

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$m_* \leftarrow \mathcal{M}$

$s_* \leftarrow \mathcal{A}^{\mathcal{O}()}(pk_*, m_*)$

if $\text{Ver}(pk_*, m_*, s_*) = 1$ **output** 1; // \mathcal{A} wins

else output 0;

Oracle $\mathcal{O}()$ // Signing a random message oracle

$m \leftarrow \mathcal{M}$

return $(m, \text{Sig}(sk_*, m))$

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$s_* \leftarrow \mathcal{A}^{\mathcal{O}()}(pk_*, m_*)$

if $\text{Ver}(pk_*, m_*, s_*) = 1$ **output** 1; // \mathcal{A} wins

else output 0;

Oracle $\mathcal{O}()$ // Signing a random message oracle

$m \leftarrow \mathcal{M}$

return $(m, \text{Sig}(sk_*, m))$

Definition (Sig-UF-RMA)

The signature scheme Σ is UF-RMA secure if for all PPTM \mathcal{A}

$$\Pr[\text{Exp-UF-RMA}(\Sigma, \mathcal{A}, \ell) = 1] \in \text{negl}(\ell)$$

UF-RMA Security

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Improvement: Make the adversary choose the target message to be signed.

Existential Unforgeability (EF)

Experiment $\text{Exp-EF-RMA}(\Sigma, \mathcal{A}, \ell)$:

$(pk_*, sk_*) \leftarrow \text{KeyGen}(\ell)$

$L = \emptyset$

$(m_*, s_*) \leftarrow \mathcal{A}^{O()}(pk_*)$

if $m_* \notin L$ **and** $\text{Ver}(pk_*, m_*, s_*) = 1$ **output** 1; // \mathcal{A} wins

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$m \leftarrow \mathcal{M}$

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We need to maintain the list of messages signed by the oracle to exclude trivial attacks.

Beyond RMA Security

Plain RSA and ElGamal signatures are EF-RMA insecure, while RSA-FDH, Pointcheval-Stern, DSA and ECDSA are EF-RMA secure.

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Improvement: Allow the adversary choose the messages to be signed by the oracle.

EF-CMA Security

Experiment $\text{Exp-EF-CMA}(\Sigma, \mathcal{A}, \ell)$:

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EF-CMA Security

EF-CMA is nowadays considered the standard security level for a digital signature scheme.

RSA-FDH, Pointcheval-Stern, DSA and ECDSA achieve EF-CMA security in the Random Oracle Model.

There are other signature schemes (e.g., Cramer-Shoup) that also achieve EF-CMA security without random oracles.

Outline

- 1 Defining Computational Security
- 2 Security Models for Public Key Encryption
- 3 Security Models for Digital Signatures
- 4 Security Assumptions and Results**

Assumptions Related to Integer Factorization

- **IF:** For random ℓ -bit long primes p, q , **given** $n = pq$, **compute** p or q .
- **RSA:** For random ℓ -bit long primes p, q , $n = pq$, a random e coprime with $\phi(n)$ and a random $m \in \mathbb{Z}_n$, **given** n and $m^e \bmod n$, **compute** m .

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- **e-RSA:** For random ℓ -bit long primes p, q such that e does not divide $p - 1$ or $q - 1$ and a random $m \in \mathbb{Z}_n$, **given** n and $m^e \bmod n$, **compute** m .

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- **DCR:** For random ℓ -bit long primes p, q , $n = pq$ and a random $r \in \mathbb{Z}_{n^2}$, **tell apart** the two probability distributions (n, r) and $(n, r^n \bmod n^2)$.

The last problem is called a **decision** problem, while the other three are **search** problems.

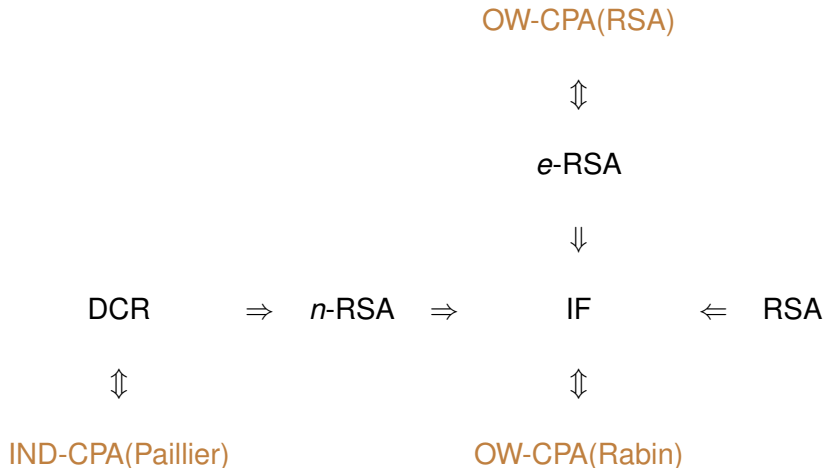
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All these problems are conjectured hard.

Known Hardness Implications



Assumptions Related to Discrete Logarithm

For a ℓ -bit long prime q , a cyclic group G of order q and a generator g :

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All these problems are conjectured hard on suitable groups like:

- Subgroups of the multiplicative group of a large enough finite field.
- Subgroups of an elliptic curve over a large enough finite field.

Known Hardness Implications

DDH \Rightarrow CDH \Rightarrow DLOG

\Updownarrow

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IND-CPA(EIGamal) \Rightarrow OW-CPA(EIGamal)

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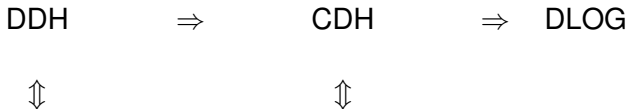
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In some special groups (e.g., elliptic curves with efficient pairing maps) DDH is easy while CDH is still conjectured to be hard.

Secure variants of ElGamal encryption and new encryption and signature schemes with new features have been proposed in these special groups (pairing based cryptography).

Further Assumptions

There exist other settings that offer well-known problems also conjectured hard, that can be used to build cryptographic protocols:

- Lattices
- Coding Theory
- Multivariate Polynomials
- Isogenies
- Non-abelian Groups

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- Lattices
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- Isogenies
- Non-abelian Groups

Some of the problems could remain hard even in the presence of quantum computers.

Data Protection

Jorge L. Villar

MCYBERS, UPC, Fall 2025

END