

On the metric dimension of graphs

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RESOLVING SETS

★ $G = (V, E)$ is a connected graph.

▶ A vertex x of G RESOLVES a pair of vertices u, v of G if:

$$d(x, u) \neq d(x, v).$$

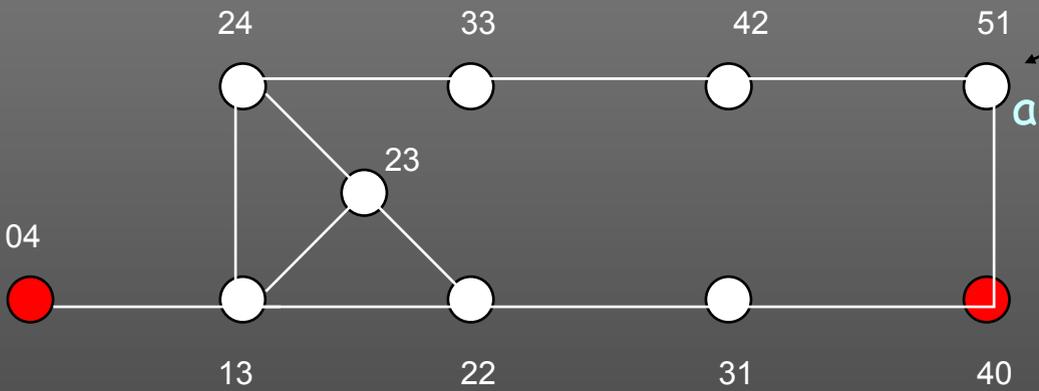
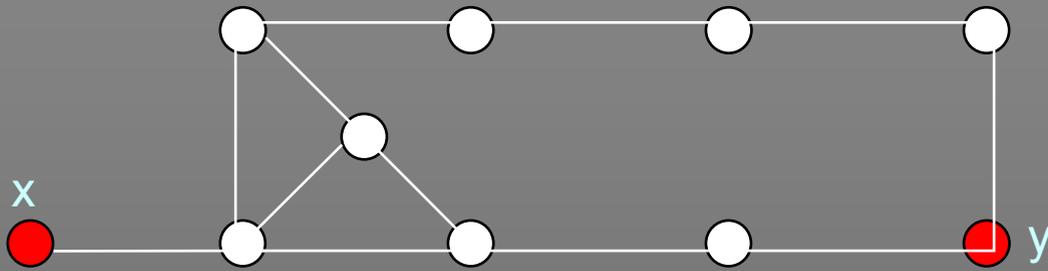
▶ A vertex subset $S \subset V$ is a RESOLVING SET of G if:

every two distinct vertices of G are resolved by some vertex of S .

METRIC DIMENSION

- ★ $S \subset V$ is a resolving set of $G = (V, E)$.
- ▶ S is a metric basis if it is a minimum resolving set.
- ★ $S = \{u_1, u_2, \dots, u_r\}$ is a metric basis of G .
- ▶ The metric dimension of G is: $\beta(G) = r$.
- ★ x is a vertex of G .
- ▶ The metric coordinates of x are:

$$(x_1, \dots, x_r), \text{ where } x_i = d(u_i, x).$$



$d(x, a) = 5$
 $d(y, a) = 1$

$\beta(G) = 2$

$S = \{x, y\}$ is a metric basis

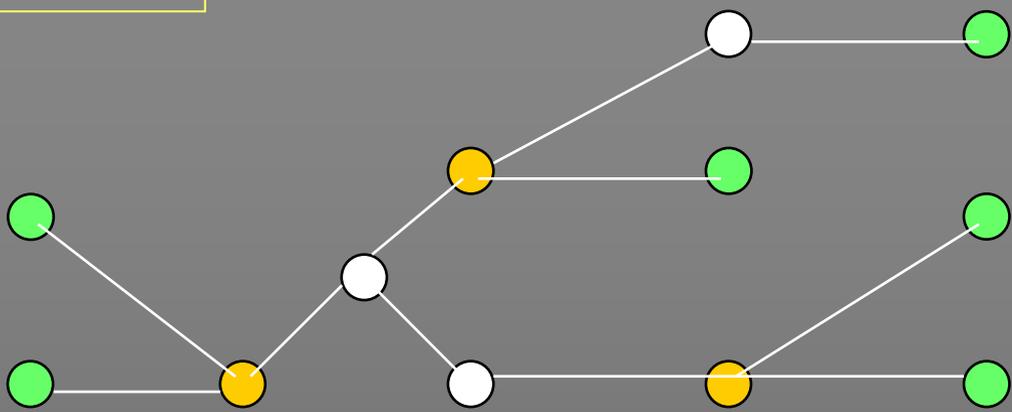
SOME BASIC KNOWN RESULTS

- The problem of computing the metric dimension of an arbitrary graph is NP-hard..

name	path	cycle	complete	bicomplete	wheel	hypercube
G	P_n	C_n	K_n	$K_{r,s}$	$W_{1,r}$	$Q_r = [K_2]^r$
$ V(G) $	$n \geq 1$	$n \geq 3$	$n \geq 2$	$r + s \geq 3$	$r + 1 = 4, 7$	2^r
$\beta(G)$	1	2	$n - 1$	$n - 2$	3	$r \ (r \leq 4)$

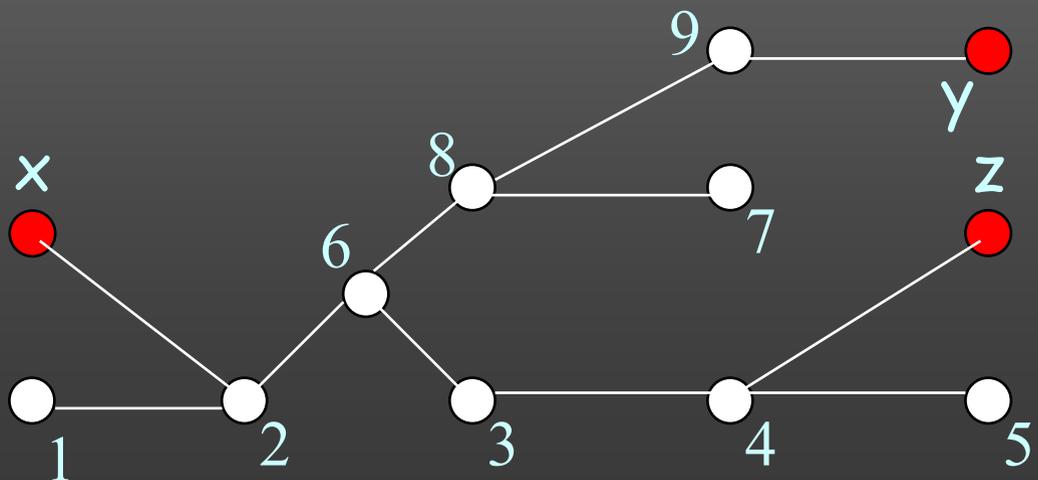
- If T is a tree s.t. $\lambda(T) \geq 1$, then $\beta(T) = |Ext(T)| - \lambda(T)$.
- If $G = T + e$, then $\beta(T) - 2 \leq \beta(T + e) \leq \beta(T) + 1$.
- $r \notin \{3, 6\}$: $\beta(W_{1,r}) = \lfloor \frac{2r+2}{5} \rfloor$.

OTUS, INTERLU



● leaves

● exterior major vertices



1	255
2	144
3	342
4	451
5	562
6	233
7	435
8	324
9	415
x	055
y	506
z	560

BOUNDARY VERTICES

★ u, v are two vertices of a graph G .

▶ v is a boundary vertex of u if:

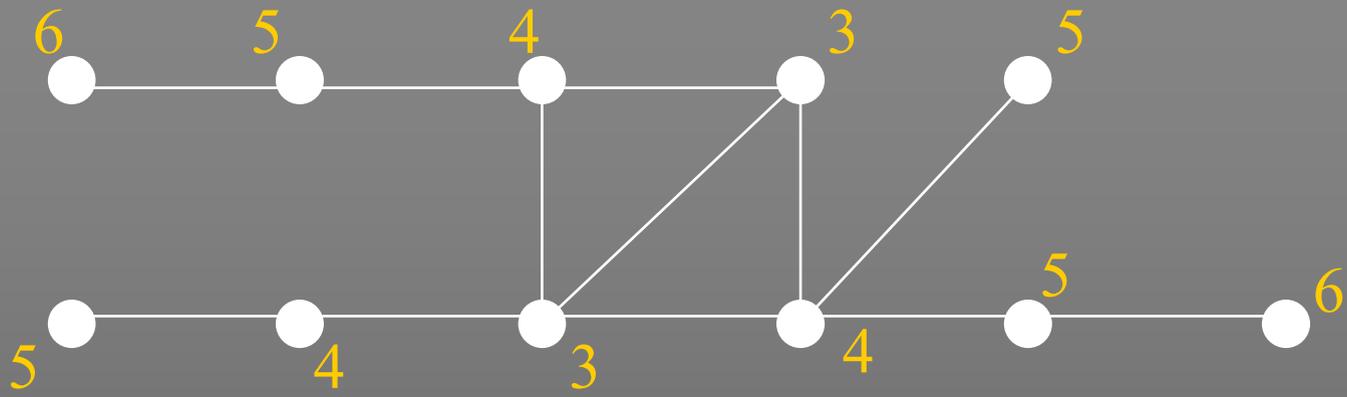
$$\forall w \in N(v), d(u, w) \leq d(u, v).$$

▶ The set of all boundary vertices of u is its boundary $\partial(u)$.

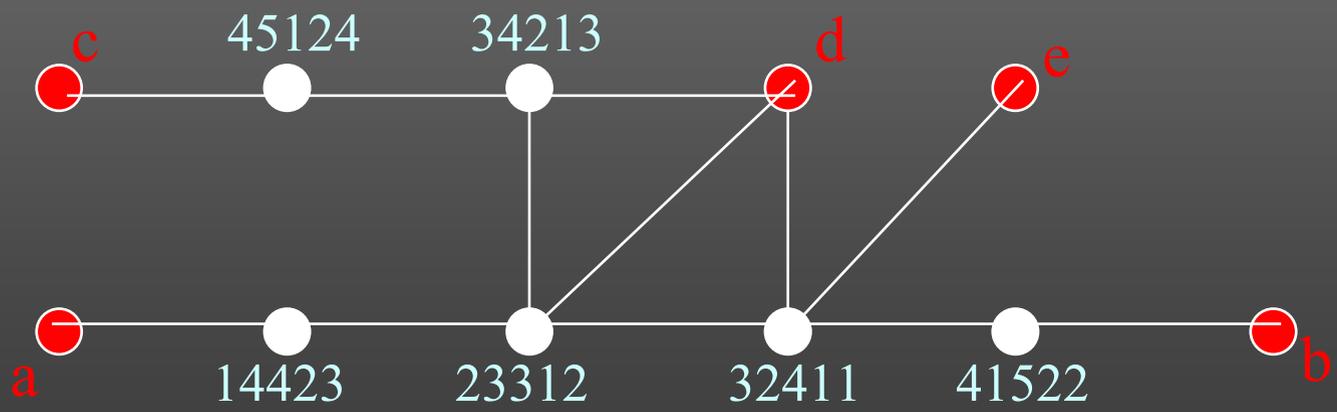
▶ The boundary of G is the set $\partial(G) = \cup_{u \in V(G)} \partial(u)$.

■ **THEOREM:** The boundary $\partial(G)$ is a resolving set of G .

■ **COROLLARY:** $\beta(G) \leq |\partial(G)|$.



$$\partial(G) = \{a, b, c, d, e\}$$



BOUNDS AND CHARACTERIZATIONS

- $|G| - d^{\beta(G)} \leq \beta(G) \leq |G| - d$ ($d = \text{diam}(G)$).
- $\beta(G) = 1 \Leftrightarrow G = P_n$.
- $\beta(G) = |G| - 1 \Leftrightarrow G = K_n$.
- $t \geq 2$: $\beta(G) = |G| - 2 \Leftrightarrow G \in \{K_{s,t}, K_s + \overline{K}_t, K_s + (K_1 \cup K_t)\}$.
- $d \geq 3$, $\beta(G) = |G| - d \Rightarrow \exists a \in \text{Per}(G)$ without twins.
- $\beta(G) = |G| - 3 \Leftrightarrow G \in \{\dots, \dots, \dots\}$.



CARTESIAN PRODUCT

★ $G_1 = (V_1, E_1)$ $G_2 = (V_2, E_2)$ are two connected graphs.

▶ The cartesian product $G_1 \square G_2$ is the graph with:

◇ $V(G_1 \square G_2) = V_1 \times V_2$.

◇ (u_1, u_2) is adjacent to (v_1, v_2) iff $\begin{cases} u_1 = v_1 \text{ and } u_2 v_2 \in E_2 \\ \text{or} \\ u_1 v_1 \in E_1 \text{ and } u_2 = v_2 \end{cases}$

■ $d((u_1, u_2), (v_1, v_2)) = d_{G_1}(u_1, v_1) + d_{G_2}(u_2, v_2)$.

■ THEOREM: $\beta(G_1 \square G_2) \geq \max[\beta(G_1), \beta(G_2)]$.

SOME KNOWN RESULTS ON THE C.P.

■ $\beta(G) \leq \beta(G \square K_2) \leq \beta(G) + 1.$

■ $\beta(P_m \square P_n) = 2.$

■ $\lim_{n \rightarrow \infty} \beta(Q_n) \cdot \frac{\log_2(n)}{n} = 2$

■ $\beta(P_{m_1} \square P_{m_2} \square \dots \square P_{m_d}) = d$ (wrong result: DAM, 70 (1996)).

CARTESIAN PRODUCTS [BOUNDS]

- $\max[\beta(G), \beta(H)] \leq \beta(G \square H)$.
- $|G| \geq 3, |H| \geq 3$: $\beta(G \square H) \leq \min\{\beta(G) + |H|, \beta(H) + |G|\} - 2$
- $\beta(G \square K_n) \leq \max\{n - 1, 2 \cdot \beta(G)\}$. ■ $\beta(G \square P_n) \leq \beta(G) + 1$.
- $\beta(G \square C_n) \leq \beta(G) + 2$. ■ $\beta(G \square T) \leq \beta(G) + |Ext(T)| - 1$.
- For all $k \geq 1$ and $n \geq 2$ there is a k -connected graph $G_{n,k}$ for which $\beta(G_{n,k}) \leq 2k$ and $\beta(G_{n,k} \square G_{n,k}) \geq n$.
- For all $k \geq 1$ there is no function f such that: $\beta(G \square H) \leq f(\beta(G), \beta(H))$, for all k -connected graphs G and H .

CARTESIAN PRODUCTS [EXACT VALUES]

■ $m \leq n \Rightarrow \dim(K_m \square K_n) = \begin{cases} n - 1 & \text{if } 2m - 2 < n, \\ \lfloor \frac{2m+2n-2}{3} \rfloor & \text{if } 2m - 2 \geq n. \end{cases}$

■ $m \geq 4: \beta(K_m \square C_n) = \begin{cases} m, & \text{if } m = 4 \text{ and } n \text{ odd,} \\ m - 1, & \text{otherwise.} \end{cases}$

■ $\beta(C_m \square C_n) = \begin{cases} 3, & \text{if } m \text{ or } n \text{ is odd} \\ 4, & \text{otherwise.} \end{cases}$

■ $\beta(C_m \square P_n) = \begin{cases} 2, & \text{if } m \text{ odd} \\ 3, & \text{if } m \text{ even (and } n \neq 1) \end{cases}$

$G \setminus H$	K_n	C_n	P_n
K_m	$n - 1, \lfloor \frac{2m+2n-2}{3} \rfloor$	$m - 1, m$	$m - 1$
C_m		3, 4	2, 3
P_m			2

CARTESIAN PRODUCT OF A CYCLE BY A PATH

■ $m \geq 3, n \geq 2: \beta(C_m) = 2, \beta(P_n) = 1$

■ $\max[\beta(G), \beta(H)] \leq \beta(G \square H).$

■ $\beta(G \square P_n) \leq \beta(G) + 1.$

$\Rightarrow 2 = \beta(C_m) \leq \beta(C_m \square P_n) \leq \beta(C_m) + 1 = 3.$

◆ $\beta(C_m \square G) = 2 \Leftrightarrow G$ is a path and m is odd.

THEOREM: $\beta(C_m \square P_n) = \begin{cases} 2, & \text{if } m \text{ odd} \\ 3, & \text{if } m \text{ even (and } n \neq 1) \end{cases}$

DOUBLY RESOLVING SETS

► Two vertices v, w of $G \neq K_1$ are DOUBLY RESOLVED by a pair of vertices x, y of G if:

$$d(v, x) - d(w, x) \neq d(v, y) - d(w, y).$$

► A set $S \subseteq V$ is a DOUBLY RESOLVING SET of G if every pair of distinct vertices of G are doubly resolved by two vertices in S . $\psi(G)$ is the minimum cardinality of a doubly resolving set.

■ $1 \leq \beta(G) \leq \psi(G) \leq |G| - 1$ ■ $2 \cdot \beta(G \square G) \geq \psi(G)$

■ $\beta(G \square H) \leq \beta(G) + \psi(H) - 1$ ■ $\beta(G \square H) \leq \beta(G) + 2\beta(H \square H) - 1$

name	path	odd cycle	even cycle	complete	tree
G	P_n	C_n	C_n	K_n	T_n
$ V(G) $	$n \geq 2$	$n \geq 3$	$n \geq 4$	$n \geq 3$	$n \geq 2$
$\beta(G)$	1	2	2	$n - 1$	$ Ext(T_n) - \lambda(T_n)$
$\psi(G)$	2	2	3	$n - 1$	$ Ext(T_n) $

C. P. OF AN ODD CYCLE BY A PATH

$$\blacksquare \psi(C_m) = \begin{cases} 2, & \text{if } m \text{ odd} \\ 3, & \text{if } m \text{ even} \end{cases}$$

$$\blacksquare \psi(P_n) = 2$$

$$\blacksquare \beta(G) \leq \beta(G \square P_n) \leq \beta(G) + \psi(P_n) - 1 = \beta(G) + 1$$

$$\blacksquare \beta(G) \leq \beta(C_m \square G) \leq \beta(G) + \psi(C_m) - 1 \leq \begin{cases} \beta(G) + 1, & \text{if } m \text{ odd} \\ \beta(G) + 2, & \text{if } m \text{ even} \end{cases}$$

$$\Rightarrow m \text{ odd} \Rightarrow \beta(C_m \square P_n) = 2.$$

<http://es.arXiv.org/find/math/1/PELAYO/0/1/0/past/3/0>